## Problem

Let  $n \in \mathbb{N}$  be an even number, say n = 2m. Also let  $\mathbb{I}_n = \{1, 2, 3, \dots, n\}$ . Let  $f : \mathbb{I}_n \to \mathbb{I}_n$  be a bijection; in other words, it represents a permutation of length n. If  $i \in \mathbb{I}_n$  such that f(i) = i, then i is called a *fixed point* of f. If  $i \in \mathbb{I}_n$  such that f(i) = n + 1 - i, then i is called a *mirror point* of f.

Let  $X_n$  denote the number of permutations of length n with no fixed points and no mirror points. Find the value of  $X_n$ .

## Solution

We wish to apply the Principle of Inclusion-Exclusion (PIE).

Let our universal set be  $S_n$ , the set of all permutations of length n; hence  $|S_n| = n!$ . For all  $i \in \mathbb{I}_n$ , we define  $A_i = \{f \in S_n | f(i) = i\}$ , and  $B_i = \{f \in S_n | f(i) = n + 1 - i\}$ . In other words,  $A_i$  is the set of all permutations for which i is a fixed point, and  $B_i$  is the set of all permutations for which i is a fixed point, and  $B_i$  is the set of all permutations for which i is a mirror point.

We note the following lemma, which will help us to make suitable cases later on: Lemma: For any  $i, j \in \mathbb{I}_n$  such that i + j = n + 1,  $|A_i \cap B_i| = |A_i \cap B_j| = 0$ .

**Proof**: It suffices to note that n is even, hence no element can be a fixed point as well as a mirror point of the same permutation.

(If n is odd, then only the middle element  $\frac{n+1}{2}$  can be a fixed point as well as a mirror point. The reader is encouraged to suitably modify this solution to work for odd values of n as well.)

By PIE, what we want to calculate is:

$$X_n = |S| - |A_1 \cup A_2 \cup \dots \cup A_n \cup B_1 \cup B_2 \cup \dots \cup B_n| = |S| + \sum_{t=1}^{2n} (-1)^t E_t$$
(1)

where  $E_t$  denotes the sum of the sizes of all t-set intersections from the  $A_i, B_j$  families. In other words,  $E_t = \sum |A_{i_1} \cap A_{i_2} \cap \cdots \cap A_{i_k} \cap B_{j_1} \cap B_{j_2} \cap \cdots \cap B_{j_l}|$ , where the summation is taken over all  $1 \leq i_1 < i_2 < \cdots < i_k \leq n$  and  $1 \leq j_1 < j_2 < \cdots < j_l \leq n$  such that k + l = t; including the possibility that either of k, l could be zero.

To calculate  $E_t$ , all such *t*-set intersections  $I = |A_{i_1} \cap A_{i_2} \cap \cdots \cap A_{i_k} \cap B_{j_1} \cap B_{j_2} \cap \cdots \cap B_{j_l}|$  can be partitioned into the following cases:

**Case 1**: There exist some  $1 \le r \le k$  and  $1 \le s \le l$ , such that  $i_r = j_s$ , or  $i_r + j_s = n + 1$ Due to the above lemma, we note that  $|A_{i_r} \cap B_{j_s}| = 0$ ; hence |I| = 0.

**Case 2**: There do not exist any  $1 \le r \le k$  and  $1 \le s \le l$ , for which  $i_r = j_s$  or  $i_r + j_s = n + 1$ . For any permutation f that belongs to the *t*-set intersection I as defined above, the following values of f are forced:

 $f(i_r) = i_r$  for all r from 1 to k, and  $f(j_s) = n + 1 - j_s$  for all s from 1 to l.

All the above values in the domain and codomain are distinct, which is consistent with f being a bijection.

All the remaining n-k-l=n-t values of f can be chosen in (n-t)! ways; so |I|=(n-t)!.

Now we only need to count the *t*-set intersections that belong to case 2.

Noting that n = 2m, we partition the 2n sets of  $A_i, B_j$  families into m groups of 4 sets each:  $\{A_1, A_n, B_1, B_n\}, \{A_2, A_{n-1}, B_2, B_{n-1}\}, \cdots$ 

 $\cdots, \{A_i, A_{n+1-i}, B_i, B_{n+1-i}\}, \cdots, \{A_m, A_{m+1}, B_m, B_{m+1}\}$ 

As per the constraint imposed by case 2, no two  $A_i, B_j$  from the same group can be used in any given *t*-set intersection. So we can use at most 2 out of the 4 elements from each group.

Let us count the sub-case in which we use exactly 2 elements from some r groups, and exactly 1 element from some s groups; where 2r + s = t.

We can choose the first r groups in  $\binom{m}{r}$  ways, and the next s groups in  $\binom{m-r}{s}$  ways.

For each of the first r groups, we can use either both elements of the  $A_i$  family, or both elements of the  $B_i$  family, in that group.

For each of the next s groups, we can use any one of the 4 elements in that group. Thus the number of t-set intersections in this sub-case is  $\binom{m}{r}\binom{m-r}{s}2^r4^s = \binom{m}{r}\binom{m-r}{t-2r}2^{2t-3r}$ .

Summing up the above sub-cases over all valid values of r, and noting that each such t-set intersection is of size (n-t)!, we get:

$$E_{t} = (n-t)! \sum_{r=0}^{\lfloor \frac{t}{2} \rfloor} {m \choose r} {m-r \choose t-2r} 2^{2t-3r}$$
(2)

We note that if the two binomial coefficients on the RHS are non-zero, then  $m \ge r$  and  $m-r \ge t-2r$ ; which forces  $t \le 2m = n$ . In other words,  $E_t = 0$  for all t > n; which makes sense because we cannot force more than n constraints on a function having n inputs. The above formula also gives  $E_0 = n!$  which is the size of our universal set  $S_n$ .

We can now combine (1) and (2) to write the formula for  $X_n$ , which is the number of *n*-length permutations with no fixed points and no mirror points:

$$X_n = \sum_{t=0}^n (-1)^t (n-t)! \sum_{r=0}^{\lfloor \frac{t}{2} \rfloor} {m \choose r} {m-r \choose t-2r} 2^{2t-3r}$$

By exchanging the order of summation, and other simplifications, we can rewrite this as:

$$X_{2m} = \sum_{s=0}^{m} \binom{m}{s} 2^{m-s} F_1(s)$$

where  $F_1(s) = \sum_{k=0}^{s} (2s-k)! \binom{s}{k} (-4)^k$ .

Similarly, the reader is encouraged to derive the following formula for the case when n is odd, say n = 2m + 1:

$$X_{2m+1} = \sum_{s=0}^{m} \binom{m}{s} 2^{m-s} (F_2(s) - F_1(s))$$

where  $F_1(s)$  is the same as above, and  $F_2(s) = \sum_{k=0}^{s} (2s - k + 1)! \binom{s}{k} (-4)^k$ .

It seems difficult to simplify these any further; the reader is encouraged to try the same.

Finally, we provide some initial values of  $X_n$ , along with the number of derangements  $D_n$ , as well as a sample implementation in Python to count  $D_n, X_n$  directly.

n	$D_n$	$X_n$
1	0	0
2	1	0
3	2	0
4	9	4
5	44	16
6	265	80
7	1854	672
8	14833	4752
9	133496	48768
10	1334961	440192

```
def get_permutations(n, isValid):
 # use backtracking to generate all permutations recursively
 def perms_step(n, output, available, isValid):
    # output: permutation generated so far
    # available: elements available to append in this step
    # isValid is a function that takes inputs as n, i, p(i),
    # and returns whether it is ok to put element p(i) at position i
    if len(available) == 0:
     return [output]
    all_outputs = []
    for p_i in available:
      if isValid(n, len(output), p_i):
        newOutput = output.copy(); newOutput.append(p_i)
        newAvailable = available.copy(); newAvailable.remove(p_i)
        all_outputs += perms_step(n, newOutput, newAvailable, isValid)
    return all_outputs
  return perms_step(n, [], set(range(n)), isValid)
def isNotFixedPt(n, i, p_i):
 return p_i != i
def isNotMirrorPt(n, i, p_i):
 return p_i != n - 1 - i
def isNotFixedOrMirrorPt(n, i, p_i):
 return isNotFixedPt(n, i, p_i) and isNotMirrorPt(n, i, p_i)
def print_dn_xn_table(N):
 for n in range(1, N):
    D_n = len(get_permutations(n, isNotFixedPt))
    X_n = len(get_permutations(n, isNotFixedOrMirrorPt))
    print(f'n:{n: 3} D_n:{D_n: 9} X_n:{X_n: 9}')
if __name__ == '__main__':
 print_dn_xn_table(11)
```

## mirror\_pts.py